Generalised Buchberger and Schreyer algorithms for strongly discrete coherent rings

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Abstract

Let M be a finitely generated submodule of a free module over a multivariate polynomial ring with coefficients in a discrete coherent ring. We prove that its module MLT(M) of leading terms is countably generated and provide an algorithm for computing explicitly a generating set. This result is also useful when MLT(M) is not finitely generated.

Suppose that the base ring is strongly discrete coherent. We provide a Buchberger-like algorithm and prove that it converges if, and only if, the module of leading terms is finitely generated. We also provide a constructive version of Hilbert's syzygy theorem by following Schreyer's method.

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1 Introduction

This paper is written in Bishop's style of constructive mathematics (see Bishop 1967, Bishop and Bridges 1985, Mines, Richman, and Ruitenburg 1988, Lombardi and Quitté 2015, Yengui 2015). It can be seen as a sequel to the paper Gamanda, Lombardi, Neuwirth, and Yengui 2020 (see also Hadj Kacem and Yengui 2010, Yengui 2006). It generalises results in Adams and Loustaunau 1994 to suitable nonnoetherian contexts.

Our general context is the following.

General context 1.1. In this article, **R** is a commutative ring with unit, X_1, \ldots, X_n are *n* indeterminates $(n \ge 1)$, $\mathbf{R}[\underline{X}] = \mathbf{R}[X_1, \ldots, X_n]$, $\mathbf{H}_n^m = \mathbf{R}[\underline{X}]^m$ is a free $\mathbf{R}[\underline{X}]$ -module with basis (e_1, \ldots, e_m) $(m \ge 1)$, and > is a monomial order on \mathbf{H}_n^m (see Definition 3.1). The case of an ideal in $\mathbf{R}[\underline{X}]$ is addressed by considering m = 1: $\mathbf{H}_n^1 = \mathbf{R}[\underline{X}]$ and e_1 is the unit of \mathbf{R} .

We shall need to qualify this context by suitable hypotheses of coherence and discreteness (see Definition 2.1). Let us consider a finitely generated submodule M of \mathbf{H}_n^m and its module MLT(M) of leading terms with respect to our monomial order.

Our first Fundamental theorem 4.11 states that MLT(M) is countably generated and provides an algorithm for computing explicitly a generating set. This result is also useful when MLT(M) is not finitely generated.

The second Fundamental theorem 6.1 indicates a precise context in which the (generalised) Buchberger criterion applies and Buchberger's algorithm computes a Gröbner basis for M.

The third Fundamental theorem 7.4 indicates a precise context in which the (generalised) Schreyer method computes a finite free resolution for M.

These are central problems since polynomials or vectors of polynomials always admit nontrivial algebraic relations and such computations are a basic step, for instance, for computing resolutions of modules.

The main results of the paper are two generalisations of classical theorems/algorithms, Buchberger's algorithm for the computation of a Gröbner basis and Schreyer's algorithm for the computation of a Gröbner bases of the first syzygy module. The generalisations are given for so-called strongly-discrete coherent rings, namely rings with membership test for finitely generated ideals and with finitely generated syzygy modules.

An important aspect of extending classical theorems in the theory of Gröbner bases from polynomial rings over fields to polynomial rings over more general algebraic structures is the emergence of more elementary arguments.

2 Constructive definitions and contexts

We start with recalling the following constructive definitions.

Definition 2.1.

- The ring \mathbf{R} is *discrete* if it is equipped with a zero test: equality is decidable.
- Let U be an **R**-module. The syzygy module of an n-tuple $(v_1, \ldots, v_n) \in U^n$ is

$$Syz(v_1, ..., v_n) := \{ (b_1, ..., b_n) \in \mathbf{R}^{1 \times n} ; b_1v_1 + \dots + b_nv_n = 0 \}.$$

The syzygy module of a 1-tuple v is the annihilator Ann(v) of v.

• An **R**-module U is *coherent* if the syzygy module of every *n*-tuple of elements of U is finitely generated, i.e. if there is an algorithm providing a finite system of generators for the syzygies, and an algorithm that represents each syzygy as a linear combination of the generators. The ring **R** is *coherent* if it is coherent as an **R**module. It is well-known that a module is coherent if, and only if, on the one hand any intersection of two finitely generated submodules is finitely generated, and on the other hand the annihilator of every element is a finitely generated ideal.

• A ring is strongly discrete if it is equipped with a membership test for finitely generated ideals, i.e. if, given $a, b_1, \ldots, b_n \in \mathbf{R}$, one can answer the question $a \in ?$ $\langle b_1, \ldots, b_n \rangle$ and, in the case of a positive answer, one can explicitly provide $c_1, \ldots, c_n \in \mathbf{R}$ such that $a = b_1c_1 + \cdots + b_nc_n$.

• **R** is a *Bézout ring* if every finitely generated ideal is principal, i.e. of the form $\langle a \rangle = \mathbf{R}a$ with $a \in \mathbf{R}$. A Bézout ring is strongly discrete if, and only if, it is equipped with a divisibility test; it is coherent if, and only if, the annihilator of any element is principal. To be a valuation ring (in the Kaplansky sense) is to be a Bézout local ring (see Lombardi and Quitté 2015, Lemma IV-7.1).

• A Bézout ring **R** is *strict* if for all $b_1, b_2 \in \mathbf{R}$ we can find $d, b'_1, b'_2, c_1, c_2 \in \mathbf{R}$ such that $b_1 = db'_1, b_2 = db'_2$, and $c_1b'_1 + c_2b'_2 = 1$. Valuation rings and Bézout domains are strict Bézout rings; a quotient or a localisation of a strict Bézout ring is again a

strict Bézout ring (see Lombardi and Quitté 2015, Exercise IV-7 pp. 220–221, solution pp. 227–228). A zero-dimensional Bézout ring is strict (because it is a "Smith ring", see Díaz-Toca, Lombardi, and Quitté 2014, Exercise XVI-9 p. 355, solution p. 526, and Lombardi and Quitté 2015, Exercise IV-8 pp. 221-222, solution p. 228).

We shall consider three contexts that are more specific than General context 1.1.

Discrete coherent context 2.2. General context 1.1 with R discrete and coherent.

Simple division context 2.3. General context 1.1 with R strongly discrete.

Remark 2.4. We could have considered a "Division with remainder context" in which we would assume moreover that we have a partial preorder $\leq_{\mathbf{R}}$ on elements of \mathbf{R} (with $a \leq_{\mathbf{R}} b$ if a = b) and a generalised division algorithm Rem for \mathbf{R} which computes, for given $c, c_1, \ldots, c_k \in \mathbf{R}$, a remainder $r_0 = c - a_1c_1 - \cdots - a_kc_k$ satisfying $r_0 = 0$ if, and only if, $c \in \langle c_1, \ldots, c_k \rangle$, and $r_0, a_1c_1, \ldots, a_kc_k \leq_{\mathbf{R}} c$ otherwise. Simple division context 2.3 might be seen as the particular case where $\leq_{\mathbf{R}}$ is equality and Rem returns cif $c \notin \langle c_1, \ldots, c_k \rangle$.

Strongly discrete coherent context 2.5. General context 1.1 with R strongly discrete and coherent.

3 Gröbner bases for modules over a discrete ring

Definition 3.1 (Monomial orders on finite-rank free $\mathbf{R}[\underline{X}]$ -modules, see Adams and Loustaunau 1994, Cox, Little, and O'Shea 2005, Yengui 2021b, General context 1.1).

(1) Monomials, terms.

• A monomial in \mathbf{H}_n^m is a vector of the form $M = \underline{X}^{\alpha} e_i$ $(1 \leq i \leq m)$, where $\underline{X}^{\alpha} = X_1^{\alpha_1} \cdots X_n^{\alpha_n}$ is a monomial in $\mathbf{R}[\underline{X}]$; the index *i* is the *position* of the monomial. The set of monomials in \mathbf{H}_n^m is denoted by \mathbb{M}_n^m , with $\mathbb{M}_n^1 \cong \mathbb{M}_n$ (the set of monomials in $\mathbf{R}[\underline{X}]$). For example, $X_1 X_2^3 e_2$ is a monomial in \mathbf{H}_n^m , but $2X_1 e_3$, $(X_1 + X_2^3) e_2$, and $X_1 e_2 + X_1 e_3$ are not.

• If $M = \underline{X}^{\alpha} e_i$ and $N = \underline{X}^{\beta} e_j$, we say that M divides N if i = j and \underline{X}^{α} divides \underline{X}^{β} . For example, $X_1 e_1$ divides $X_1 X_2 e_1$, but does not divide $X_1 X_2 e_2$. Note that in the case that M divides N, there exists a monomial \underline{X}^{γ} in \mathbb{M}_n such that $N = \underline{X}^{\gamma} M$: in this case we define $N/M \coloneqq \underline{X}^{\gamma}$; for example, $(X_1 X_2 e_1)/(X_1 e_1) = X_2$.

• A term in \mathbf{H}_n^m is a vector of the form cM, where $c \in \mathbf{R} \setminus \{0\}$ and $M \in \mathbb{M}_n^m$. We say that a term cM divides a term c'M', with $c, c' \in \mathbf{R} \setminus \{0\}$ and $M, M' \in \mathbb{M}_n^m$, if c divides c' and M divides M'.

(2) A monomial order on \mathbf{H}_n^m is a relation > on \mathbb{M}_n^m such that

- > is a total order on \mathbb{M}_n^m ;
- $\underline{X}^{\alpha}M > M$ for all $M \in \mathbb{M}_n^m$ and $\underline{X}^{\alpha} \in \mathbb{M}_n \setminus \{1\};$
- $M > N \implies \underline{X}^{\alpha}M > \underline{X}^{\alpha}N$ for all $M, N \in \mathbb{M}_n^m$ and $\underline{X}^{\alpha} \in \mathbb{M}_n$.

Note that, when specialised to the case m = 1, this definition coincides with the definition of a monomial order on $\mathbf{R}[\underline{X}]$.

(3) Let the ring **R** be discrete. Any nonzero vector $u \in \mathbf{H}_n^m$ can be written as a sum of terms

$$u = c_t M_t + c_{t-1} M_{t-1} + \dots + c_1 M_1$$

with $c_1, ..., c_t \in \mathbf{R} \setminus \{0\}, M_1, ..., M_t \in \mathbb{M}_n^m$, and $M_t > M_{t-1} > \cdots > M_1$.

• We define the *leading coefficient*, *leading monomial*, and *leading term* of u as in the ring case: $LC(u) = c_t$, $LM(u) = M_t$, LT(u) = LC(u) LM(u).

• Letting $M_t = \underline{X}^{\alpha} e_{\ell}$ with $\underline{X}^{\alpha} \in \mathbb{M}_n$ and $1 \leq \ell \leq m$, we say that α is the *multidegree of* u and write $\operatorname{mdeg}(u) = \alpha$, and that the index ℓ is the *leading position* of u, and write $\operatorname{LPos}(u) = \ell$.

• We stipulate that LC(0) = 0, LM(0) = 0, and $mdeg(0) = -\infty$, but we do not define LPos(0).

(4) A monomial order on $\mathbf{R}[\underline{X}]$ gives rise to the two following canonical monomial orders on \mathbf{H}_n^m . Let us consider monomials $M = \underline{X}^{\alpha} e_i$ and $N = \underline{X}^{\beta} e_j \in \mathbb{M}_n^m$.

• We say that

$$M >_{\text{TOP}} N$$
 if either $\underline{X}^{\alpha} > \underline{X}^{\beta}$
or both $\underline{X}^{\alpha} = \underline{X}^{\beta}$ and $i < j$.

This monomial order is called *term over position* (TOP) because it gives precedence to the monomial order on $\mathbf{R}[\underline{X}]$ over the monomial position. For example, when $X_2 > X_1$, we have

$$X_2e_1 >_{\text{TOP}} X_2e_2 >_{\text{TOP}} X_1e_1 >_{\text{TOP}} X_1e_2.$$

• We say that

$$M >_{\text{POT}} N$$
 if either $i < j$
or both $i = j$ and $\underline{X}^{\alpha} > \underline{X}^{\beta}$.

This monomial order is called *position over term* (POT) because it gives precedence to the monomial position over the monomial order on $\mathbf{R}[\underline{X}]$. For example, when $X_2 > X_1$, we have

$$X_2 e_1 >_{\text{POT}} X_1 e_1 >_{\text{POT}} X_2 e_2 >_{\text{POT}} X_1 e_2$$

Definition 3.2 (list and module of leading terms, Gröbner bases). Let **R** be a discrete ring and consider a list $G = g_1, \ldots, g_p$ in \mathbf{H}_n^m . We denote by $\mathrm{LT}(G) = \mathrm{LT}(g_1), \ldots, \mathrm{LT}(g_p)$ the list of its leading terms. Suppose now that the g_i 's are nonzero and consider the finitely generated submodule $U = \langle G \rangle = \mathbf{R}[\underline{X}]g_1 + \cdots + \mathbf{R}[\underline{X}]g_p$ of \mathbf{H}_n^m .

(1) The module of leading terms of U is $MLT(U) := \langle LT(u) ; u \in U \rangle$.

(2) G is a Gröbner basis for U if $MLT(U) = \langle LT(G) \rangle$.

The following proposition comes from Gamanda, Lombardi, Neuwirth, and Yengui 2020. We give it as a motivating example showing that the "obvious" syzygies do not suffice to generate all of them.

Proposition 3.3. Let **R** be a strict Bézout ring, and $a_1, \ldots, a_s \in \mathbf{R} \setminus \{0\}$. Denote by $(\epsilon_1, \ldots, \epsilon_s)$ the canonical basis of \mathbf{R}^s . For $j \neq i$, write $a_j = d_{i,j}a_{i,j}$ with $d_{i,j} = \gcd(a_i, a_j)$. Then $\operatorname{Syz}(a_1, \ldots, a_s)$ is generated by the $\binom{s}{2}$ vectors $a_{i,j}\epsilon_i - a_{j,i}\epsilon_j$ with $i < j,^1$ together with all the $z\epsilon_i$ with $z \in \operatorname{Ann}(a_i)$. In particular, **R** is coherent if and only if $\operatorname{Ann}(a)$ is finitely generated (and thus can be generated by just one element) for any $a \in \mathbf{R}$. In that case, letting $\operatorname{Ann}(a_k) = \langle b_k \rangle$ for $1 \le k \le s$, we have:

$$\operatorname{Syz}(a_1, \ldots, a_s) = \langle a_{i,j} \epsilon_i - a_{j,i} \epsilon_j, b_k \epsilon_k ; 1 \le i < j \le s, 1 \le k \le s \rangle.$$

Proof. Let $c_1\epsilon_1 + \cdots + c_s\epsilon_s \in \text{Syz}(a_1, \ldots, a_s)$, and let $s(a_i, a_j) := a_{i,j}\epsilon_i - a_{j,i}\epsilon_j$. Note that $\text{gcd}(a_{i,j}, a_{j,i}) = 1$. For each permutation i_1, \ldots, i_s of $1, \ldots, s$, consider the product $a_{i_1,i_2} \cdots a_{i_{s-1},i_s}$. We claim that there is a Bézout identity for these products, so that it suffices to rewrite the expression $a_{i_1,i_2} \cdots a_{i_{s-1},i_s}(c_1\epsilon_1 + \cdots + c_s\epsilon_s)$ in terms of $s(a_{i_1}, a_{i_2}), \ldots, s(a_{i_{s-1}}, a_{i_s})$ and $\text{Ann}(a_{i_s})\epsilon_{i_s}$: let us replace successively

$$\begin{array}{cccc} a_{i_{1},i_{2}}\epsilon_{i_{1}} & \text{by} & \mathrm{s}(a_{i_{1}},a_{i_{2}}) + a_{i_{2},i_{1}}\epsilon_{i_{2}}, \\ & \vdots & & \vdots \\ a_{i_{s-1},i_{s}}\epsilon_{i_{s-1}} & \text{by} & \mathrm{s}(a_{i_{s-1}},a_{i_{s}}) + a_{i_{s},i_{s-1}}\epsilon_{i_{s}}. \end{array}$$

At the end, the sum will be a linear combination of $s(a_{i_1}, a_{i_2})$, $s(a_{i_2}, a_{i_3})$, ..., $s(a_{i_{s-1}}, a_{i_s})$, and ϵ_{i_s} ; let z be the coefficient of ϵ_{i_s} in this combination. As $c_1\epsilon_1 + \cdots + c_s\epsilon_s \in \text{Syz}(a_1, \ldots, a_s)$, we have $z\epsilon_{i_s} \in \text{Syz}(a_1, \ldots, a_s)$, i.e. $za_{i_s} = 0$.

It remains to obtain the Bézout identity for the products $a_{i_1,i_2} \cdots a_{i_{s-1},i_s}$. For this, it is enough to develop the product of the $\binom{s}{2}$ Bézout identities with respect to $a_{i,j}$ and $a_{j,i}$, $1 \leq i < j \leq s$: this yields a sum of products of $\binom{s}{2}$ terms, each of which is either $a_{i,j}$ or $a_{j,i}$, $1 \leq i < j \leq s$, so that it is indexed by the tournaments on the vertices $1, \ldots, s$ (i.e. the directed graphs with exactly one edge between each two vertices); every such product contains a product of the above form $a_{i_1,i_2} \cdots a_{i_{s-1},i_s}$ because every tournament contains a hamiltonian path (see Rédei 1934–1935).

¹These are the *obvious* syzygies.

Remark 3.4. The above proof results from an analysis of the following proof in the case where **R** is local, which entails in fact the general case. Since **R** is a valuation ring, we may consider a permutation i_1, \ldots, i_s of $1, \ldots, s$ such that $a_{i_s} |a_{i_{s-1}}| \cdots |a_{i_1}$. Thus $s(a_{i_1}, a_{i_2}) = \epsilon_{i_1} - a_{i_2,i_1} \epsilon_{i_2}, \ldots, s(a_{i_{s-1}}, a_{i_s}) = \epsilon_{i_{s-1}} - a_{i_s,i_{s-1}} \epsilon_{i_s}$ for some $a_{i_2,i_1}, \ldots, a_{i_s,i_{s-1}}$. Then, by replacing successively ϵ_{i_k} by $s(a_{i_k}, a_{i_{k+1}}) + a_{i_{k+1},i_k} \epsilon_{i_{k+1}}$, the syzygy (c_1, \ldots, c_s) may be rewritten as a linear combination of $s(a_{i_1}, a_{i_2}), \ldots, s(a_{i_{s-1}}, a_{i_s})$, and ϵ_{i_s} , with the coefficient of ϵ_{i_s} turning out to lie in $Ann(a_{i_s})$.

In the following, we give examples of coherent rings over which syzygy modules are not always generated by vectors with at most 2 nonzero components.

Example 3.5. Consider the (noetherian) coherent ring $\mathbb{Z}[u]$ and the syzygy module $\operatorname{Syz}(2, u, u + 2)$. As $\operatorname{Syz}(2, u) = \langle (u, -2) \rangle$, $\operatorname{Syz}(2, u + 2) = \langle (u + 2, -2) \rangle$, and $\operatorname{Syz}(u + 2, u) = \langle (-u, u + 2) \rangle$, we conclude that if $s = (s_1, s_2, s_3) \in \operatorname{Syz}(2, u, u + 2)$ can be written as a $\mathbb{Z}[u]$ -linear combination of syzygies in $\operatorname{Syz}(2, u, u + 2)$ with at most 2 nonzero components, then it has entries s_i in $\langle 2, u \rangle$. The syzygy $(1, 1, -1) \in \operatorname{Syz}(2, u, u + 2)$ does not satisfy this property.

Example 3.6. Consider the ring $\mathbf{R} = \mathbb{Z}[u] + v \mathbb{Q}(u)[v]_{(v)}$. It is coherent by Dobbs and Papick (1976, Theorem 3, since q.f.($\mathbb{Z}[u]$) = $\mathbb{Q}(u)$ and $\mathbb{Z}[u]$ is coherent) but nonnoetherian (since $\mathbb{Z}[u]$ is not a field, see Gilmer 1972, § 17, Exercise 14). As in Example 3.5, $(1, 1, -1) \in \operatorname{Syz}_{\mathbf{R}}(2, u, u + 2)$ cannot be written as an **R**-linear combination of syzygies in $\operatorname{Syz}_{\mathbf{R}}(2, u, u + 2)$ with at most 2 nonzero components (suppose so and take v = 0).

4 Syzygies in a polynomial ring over a discrete coherent ring

Definition 4.1 (syzygies of terms, Discrete coherent context 2.2). Let $p \ge 1$ and $\mathscr{P}_p = \{E : \emptyset \ne E \subseteq [1 \dots p]\}$ be the set of nonempty subsets of the set of indices $[1 \dots p] = \{1, \dots, p\}$. Consider $M_1 = M'_1 e_{i_1}, \dots, M_p = M'_p e_{i_p}$ monomials in \mathbf{H}_n^m , $a_1, \dots, a_p \in \mathbf{R}$. Let $\mathscr{P}(M_1, \dots, M_p) \subseteq \mathscr{P}_p$ be the subset of those E which are position level sets of (M_1, \dots, M_p) , i.e. such that $i_j = i_{j'}$ for $j, j' \in E$. Note that all singletons belong to $\mathscr{P}(M_1, \dots, M_p)$. For each position level set E of (M_1, \dots, M_p) , let $s_1^E, \dots, s_{\ell^E}^E$ be a finite number of generators of $\operatorname{Syz}((a_j)_{j\in E})$ as given by a certificate of coherence for \mathbf{R} ; here $s_i^E = (s_{i,j}^E)_{j\in E}$. Let $M^E = \operatorname{lcm}(M_j; j \in E)$ and $S^E(a_1M_1, \dots, a_pM_p)$ be the list $S_1^E, \dots, S_{\ell^E}^E$, where $S_i^E = (S_{i,1}^E, \dots, S_{i,p}^E)$ with

$$S_{i,j}^E = \begin{cases} s_{i,j}^E \ M^E / M_j & \text{if } j \in E, \\ 0 & \text{otherwise.} \end{cases}$$

This is a syzygy for (a_1M_1, \ldots, a_pM_p) : see Equation (1) below. Finally, let $S(a_1M_1, \ldots, a_pM_p)$ be the concatenation of all the lists $S^E(a_1M_1, \ldots, a_pM_p)$ when E ranges over the position level sets E of (M_1, \ldots, M_p) .

Example 4.2. Let $T_1 = (2X^2Y, 0) = 2M_1$, $T_2 = (XY^2, 0) = M_2$, $T_3 = (0, 4X) = 4M_3$ in $(\mathbb{Z}/8\mathbb{Z})[X, Y]^2$. We have $\mathscr{P}(M_1, M_2, M_3) = \{\{1\}, \{2\}, \{3\}, \{1, 2\}\}, S_1^{\{1\}} = (4, 0, 0), S_1^{\{2\}} = (0, 0, 0), S_1^{\{3\}} = (0, 0, 2), S_1^{\{1, 2\}} = (Y, 6X, 0), \text{ and } \ell^{\{1\}} = \ell^{\{2\}} = \ell^{\{3\}} = \ell^{\{1, 2\}} = 1.$

Definition 4.3 (leading monomial of summands and their leading monomial index set). Let $f_1, \ldots, f_p \in \mathbf{H}_n^m$ not all zero with $\mathrm{LC}(f_j) = a_j$ and $\mathrm{LM}(f_j) = M_j$. Let $g_1, \ldots, g_p \in \mathbf{R}[\underline{X}]$ with $\mathrm{LC}(g_j) = b_j$ and $\mathrm{LM}(g_j) = N_j$. The leading monomial of the summands of $g_1f_1 + \cdots + g_pf_p$ is the monomial $L = L(\underline{g}) = \sup_{j \in [1..p]} N_j M_j$,² and their leading monomial index set is $E = \{j; N_j M_j = L\}$.

Propositions 4.4 and 4.6 as well as Fundamental theorem 4.11 generalise the method of Gamanda, Lombardi, Neuwirth, and Yengui 2020, Theorem 4.5.

Proposition 4.4 (Discrete coherent context 2.2). The finite list $S(a_1M_1, \ldots, a_pM_p)$ generates the syzygy module $Syz(a_1M_1, \ldots, a_pM_p) \subseteq \mathbf{R}[\underline{X}]^p$.

Proof. Let us use the notation of Definition 4.1. We first check that each S_i^E is a syzygy for (a_1M_1, \ldots, a_pM_p) :

$$S_{i,1}^{E}a_{1}M_{1} + \dots + S_{i,p}^{E}a_{p}M_{p} = \sum_{j \in E} s_{i,j}^{E}a_{j}M^{E} = \left(\sum_{j \in E} s_{i,j}^{E}a_{j}\right)M^{E} = 0.$$
(1)

Conversely, let $\underline{g} = (g_1, \ldots, g_p) \in \text{Syz}(a_1M_1, \ldots, a_pM_p)$, not all g_j zero, let $L = L(\underline{g})$ be the leading monomial of the summands of $g_1a_1M_1 + \cdots + g_pa_pM_p$, and let \overline{E} be their leading monomial index set. We have $\sum_{j \in E} b_ja_j = 0$, and thus $(b_j)_{j \in E} = c_1s_1^E + \cdots + c_{\ell^E}s_{\ell^E}^E$ for some $c_1, \ldots, c_{\ell^E} \in \mathbf{R}$. Let

$$\underline{g'} = (g'_1, \dots, g'_p) = \underline{g} - \frac{L}{M^E} \sum_{i=1}^{\ell^E} c_i S_i^E \in \operatorname{Syz}(a_1 M_1, \dots, a_p M_p);$$

note that M^E divides L because every M_j , $j \in E$, does. We have $g'_j = g_j$ for $j \notin E$ and, for $j \in E$,

$$g'_{j} = g_{j} - \frac{L}{M^{E}} \sum_{i=1}^{\ell^{E}} c_{i} S_{i,j}^{E} = g_{j} - \frac{L}{M^{E}} \sum_{i=1}^{\ell^{E}} c_{i} \frac{M^{E}}{M_{j}} s_{i,j}^{E}$$
$$= g_{j} - \frac{L}{M_{j}} \sum_{i=1}^{\ell^{E}} c_{i} s_{i,j}^{E} = g_{j} - \frac{L}{M_{j}} b_{j} = g_{j} - \text{LT}(g_{j}).$$

Thus $L(\underline{g'}) < L(\underline{g})$. Reiterating this (with $\underline{g'}$ instead of \underline{g}), we reach the desired result after a finite number of steps since the set of monomials is well ordered.

 $^{^{2}}$ Note that this is not in general the leading monomial of the vector that is the sum of the expression.

Example 4.5 (Example 4.2 continued). In $(\mathbb{Z}/8\mathbb{Z})[X,Y]^2$, we have

$$Syz(T_1, T_2, T_3) = \langle (4, 0, 0), (0, 0, 2), (Y, 6X, 0) \rangle$$

Following in detail the first step in the preceding proof we get the following proposition.

Proposition 4.6 (notation of Definition 4.1, Discrete coherent context 2.2). Let

$$u = \sum_{j \in [1..p]} g_j f_j, \ L = \sup_{j \in [1..p]} N_j M_j, \ and \ E = \{ j \in [1..p] ; N_j M_j = L \}.$$

If LM(u) < L, then $f_{p+1}, \ldots, f_{p+\ell^E} \in \mathbf{H}_n^m$ and $g_{p+1}, \ldots, g_{p+\ell^E} \in \mathbf{R}[\underline{X}]$ defined by

$$f_{p+i} = \sum_{j \in E} S_{i,j}^E f_j \text{ and } g_{p+i} = c_i \frac{L}{M^E}$$

are such that

$$u = \sum_{j \in E} (g_j - \operatorname{LT}(g_j)) f_j + \sum_{j \in [1..p + \ell^E] \setminus E} g_j f_j$$

is an expression for u whose summands have leading monomial < L.

Proof. As LM(u) < L, the coefficient of L in $\sum_{j \in E} g_j f_j$ vanishes, so that $\sum_{j \in E} b_j a_j = 0$: we have

$$\sum_{j \in E} \operatorname{LT}(g_j) f_j = \sum_{j \in E} \sum_{1 \le i \le \ell^E} c_i s_{i,j}^E N_j f_j$$
$$= \sum_{1 \le i \le \ell^E} c_i \sum_{j \in E} S_{i,j}^E \frac{M_j N_j}{M^E} f_j = \sum_{1 \le i \le \ell^E} g_{p+i} f_{p+i}$$

with

$$\operatorname{LM}(g_{p+i})\operatorname{LM}(f_{p+i}) \le \frac{L}{M^E}\operatorname{LM}\left(\sum_{j\in E} S_{i,j}^E f_j\right) < \frac{L}{M^E}M^E = L.$$

Syzygies of terms, examples in the case of an ideal

This case is Discrete coherent context 2.2 with m = 1; every subset of \mathscr{P}_p is a position level set.

Example 4.7. Let us consider the following syzygy of $(6XY^2, 15X^2YZ, 10Z^2)$ in $\mathbb{Z}[X, Y, Z]$:

$$\underline{g} = (g_1, g_2, g_3) = (5XZ + 10Z^2, -2Y + 2Z, -3X^2Y - 6XY^2).$$

Following the algorithm given in the proof of Proposition 4.4 and considering the graded monomial lexicographic order with X > Y > Z, we have $L(\underline{g}) = L = X^2 Y^2 Z$, $E = \{1, 2\}$, $\operatorname{Syz}(6, 15) = \langle s_1^E = \frac{1}{3}(-15, 6) = (-5, 2) \rangle$, $\ell^E = 1$, $M^E = X^2 Y^2 Z$, $S_1^E = (-5\frac{X^2Y^2Z}{XY^2}, 2\frac{X^2Y^2Z}{X^2YZ}, 0) = (-5XZ, 2Y, 0)$, $(b_1, b_2) = (5, -2) = (-1) \cdot s_1^E$,

$$c_{1} = -1, \underline{g}' = (g_{1}', g_{2}', g_{3}') = \underline{g} - \frac{L}{M^{E}} \sum_{i=1}^{\ell^{E}} c_{i} S_{i}^{E} = \underline{g} + \frac{X^{2} Y^{2} Z}{X^{2} Y^{2} Z} S_{1}^{E} = \underline{g} + (-5XZ, 2Y, 0) = (10Z^{2}, 2Z, -3X^{2}Y - 6XY^{2}) = (g_{1} - \text{LT}(g_{1}), g_{2} - \text{LT}(g_{2}), g_{3}), \text{ with } L(\underline{g}') = L' = X^{2} Y Z^{2} < L(g).$$

Continuing with $\underline{g'}$, we obtain $E' = \{2,3\}$, $\operatorname{Syz}(15,10) = \langle s_1^{E'} \rangle$ with $s_1^{E'} = \frac{1}{5}(-10,15) = (-2,3), \ \ell_{E'} = 1, \ M_{E'} = X^2 Y Z^2, \ S_1^{E'} = \left(0,-2\frac{X^2 Y Z^2}{X^2 Y Z}, 3\frac{X^2 Y Z^2}{Z^2}\right) = (0,-2Z,3X^2Y), \ (b'_2,b'_3) = (2,-3) = (-1) \cdot s_1^{E'}, \ \underline{g''} = \underline{g'} - \frac{L'}{M_{E'}} \sum_{i=1}^{\ell_{E'}} c'_i S_i^{E'} = \underline{g'} + \frac{X^2 Y Z^2}{X^2 Y Z^2} S_1^{E'} = \underline{g'} + (0,-2Z,3X^2Y) = (10Z^2, 0,-6XY^2) = (g'_1,g'_2 - \operatorname{LT}(g'_2),g'_3 - \operatorname{LT}(g'_3)) = -2S_1^{\{1,3\}}.$ We conclude that

$$\underline{g} = -S_1^{\{1,2\}} - S_1^{\{2,3\}} - 2S_1^{\{1,3\}}.$$

Example 4.8. Let us consider the following syzygy of (3XY, 3Y, X) in $\mathbb{Z}[X, Y]$:

$$\underline{g} = (g_1, g_2, g_3) = (2X + Y, -3X^2 + 2XY, 3XY - 9Y^2).$$

Following the algorithm given in the proof of Proposition 4.4 and considering the lexicographic monomial order with X > Y, we have $L(\underline{g}) = L = X^2Y$, $E = \{1, 2, 3\}$, $Syz(3, 3, 1) = \langle s_1^E = (-1, 1, 0), s_2^E = (-1, 0, 3) \rangle$, $\ell^E = 2$, $M^E = XY$, $S_1^E = (-1, X, 0), S_2^E = (-1, 0, 3Y)$, $(b_1, b_2, b_3) = (2, -3, 3) = -3s_1^E + s_2^E$, $(c_1, c_2) = (-3, 1)$, $\underline{g'} = (g'_1, g'_2, g'_3) = \underline{g} - \frac{L}{M^E} \sum_{i=1}^{\ell^E} c_i S_i^E = \underline{g} - \frac{X^2Y}{XY} (-3S_1^E + S_2^E) = \underline{g} - X(2, -3X, 3Y) = (Y, 2XY, -9Y^2) = (g_1 - LT(g_1), g_2 - LT(g_2), g_3 - LT(g_3)) = 2YS_1^E - 3YS_2^E$, with $L(\underline{g'}) = XY^2 < L(\underline{g})$. We conclude that

$$\underline{g} = (-3X + 2Y)S_1^{\{1,2,3\}} + (X - 3Y)S_2^{\{1,2,3\}}$$

S-lists and iterated S-lists, a fundamental theorem

Definition 4.9 (Discrete coherent context 2.2). Let $f_1, \ldots, f_p \in \mathbf{H}_n^m \setminus \{0\}$ and consider their leading terms $a_1M_1, \ldots, a_pM_p \in \mathbf{H}_n^m$. If S_1, \ldots, S_ℓ is the list of generators of $\operatorname{Syz}(\operatorname{LT}(f_1, \ldots, f_p))$ computed in Proposition 4.4, the *S*-list of f_1, \ldots, f_p is the list

$$\mathscr{S}(f_1, \dots, f_p) = S_{1,1}f_1 + \dots + S_{1,p}f_p, \dots, S_{\ell,1}f_1 + \dots + S_{\ell,p}f_p$$

after having deleted the vanishing items. By induction, we define the *iterated S-lists* by

$$\begin{cases} \mathscr{S}^{0}(f_{1},\ldots,f_{p})=f_{1},\ldots,f_{p};\\ \mathscr{S}^{q+1}(f_{1},\ldots,f_{p})=\text{the concatenation of }\mathscr{S}^{q}(f_{1},\ldots,f_{p}) \text{ with } \mathscr{S}(\mathscr{S}^{q}(f_{1},\ldots,f_{p})).\end{cases}$$

Note that each item of an iterated S-list is in $\langle f_1, \ldots, f_p \rangle$.

Remark 4.10. If **R** is a Bézout ring then for any $a_1, \ldots, a_q \in \mathbf{R}$ there exists a finite generating set for Syz (a_1, \ldots, a_q) whose vectors have at most two nonzero components (see Proposition 3.3). Choose this generating set of syzygies in \mathbf{R}^q . It follows that in the corresponding iterated S-lists of f_1, \ldots, f_p there are only S-pairs (#E = 2) and auto-S-polynomials (#E = 1), as expected. Similarly, if **R** is a Prüfer domain (e.g. $\mathbf{R} = \{f \in \mathbb{Q}[X] ; f(\mathbb{Z}) \subseteq \mathbb{Z}\}$, which has Krull dimension equal to 2, see Lombardi 2010, Ducos 2015), then for any $a_1, \ldots, a_q \in \mathbf{R}$ there exists a finite generating set for Syz (a_1, \ldots, a_q) whose vectors have at most two nonzero components: in fact, a Prüfer domain is locally a valuation domain (thus, locally a Bézout domain). So in the corresponding iterated S-lists of f_1, \ldots, f_p there are only S-pairs (the auto-Spolynomials vanish since the ring **R** is supposed to be integral).

Fundamental theorem 4.11 (Discrete coherent context 2.2). Let $f_1, \ldots, f_p \in \mathbf{H}_n^m \setminus \{0\}$. For any $u \in \langle f_1, \ldots, f_p \rangle$ there exist $q \in \mathbb{N}$ and items p_1, \ldots, p_t in the list $\mathcal{G}^q(f_1, \ldots, f_p)$ such that $\mathrm{LT}(u) \in \langle \mathrm{LT}(p_1, \ldots, p_t) \rangle$. In other words,

$$\mathrm{MLT}(\langle f_1, \dots, f_p \rangle) = \bigcup_{q \in \mathbb{N}} \langle \mathrm{LT}(\mathscr{S}^q(f_1, \dots, f_p)) \rangle.$$

Comment 4.12. Compared to Theorem 4.2.8 of Adams and Loustaunau (1994), who suppose that the base ring \mathbf{R} is strongly discrete, coherent, and noetherian, our Theorem 4.11 supposes only that \mathbf{R} is discrete and coherent. Moreover, we do not perform divisions. This could be useful when one tries to prove results on the structure of the leading terms ideals (see Ben Amor and Yengui 2021, Guyot and Yengui 2024, Yengui 2021a, 2022 and the recent solution of the Gröbner ring conjecture in Yengui 2024). However Theorem 4.11 does not give a termination condition when one knows that the leading terms ideal is finitely generated (such a condition is given in Theorem 6.1). Our Theorem 4.11 is also useful when the leading terms ideal is not finitely generated (see Example 4.14 (1) and the counterexample given in Yengui 2021a).

Proof. Write

$$u = \sum_{j=1}^{p} g_j f_j \text{ with } N_j = \text{LM}(g_j) \text{ and } M_j = \text{LM}(f_j).$$
(2)

So $LM(u) \leq \sup_{1 \leq j \leq p} (N_j M_j) =: L$ (the leading monomial of the summands of u in (2)).

Case 1. LM(u) = L. Clearly $LT(u) \in \langle LT(f_1, \dots, f_p) \rangle$.

Case 2. LM(u) < L. Let $E = \{j ; N_jM_j = L\}$. By virtue of Proposition 4.6, we obtain another expression for u,

$$u = \sum_{j \in E} (g_j - \operatorname{LT}(g_j)) f_j + \sum_{j \in [1..p + \ell^E] \setminus E} g_j f_j$$
(3)

with the f_j 's in $\mathscr{S}^1(f_1, \ldots, f_p)$ and the g_j 's in $\mathbf{R}[\underline{X}]$, and the leading monomial of the summands of u in (3) is $\langle L$. Reiterating this, we end up with a situation like that

of Case 1 because the set of monomials is well ordered. So we reach the desired result after a finite number of steps. $\hfill \Box$

Remark 4.13. In the proof of Fundamental theorem 4.11 with m = 1, if the considered monomial order refines total degree (i.e. if M > N whenever $\operatorname{tdeg}(M) > \operatorname{tdeg}(N)$), then, letting $d = \max_{1 \le j \le p} (\operatorname{tdeg}(g_j) + \operatorname{tdeg}(f_j))$ and $\delta = \operatorname{tdeg}(u)$ (assumed ≥ 1), we have $q \le \binom{n+d}{d} - \binom{n+\delta-1}{\delta-1}$ (the number of monomials in X_1, \ldots, X_n of total degree at least δ and at most d).

Example 4.14. Let **V** be a nonarchimedean valuation domain, i.e. a valuation domain **V** such that there exist nonunits $a, b \in \mathbf{V}$ with a^q dividing b for every $q \in \mathbb{N}$.

(1) Let $f_1 = aX + 1$, $f_2 = b \in \mathbf{V}[X]$. Then $\mathrm{MLT}(\langle f_1, f_2 \rangle)$ is not finitely generated (see Yengui 2015, Example 253): $\mathrm{LT}(\mathscr{S}^q(f_1, f_2)) = \langle aX, b, \frac{b}{a}, \dots, \frac{b}{a^q} \rangle$ and $\mathrm{MLT}(\langle f_1, f_2 \rangle) = \langle aX, b, \frac{b}{a}, \frac{b}{a^2}, \dots \rangle$.

(2) Let $f_1 = a^2 + aXY$, $f_2 = bY^2 \in \mathbf{V}[X, Y]$. We have

$$\begin{split} \left\langle \mathrm{LT}(\mathscr{S}^0(f_1, f_2) \right\rangle &= aY \left\langle X, \frac{b}{a}Y \right\rangle \subsetneq \left\langle \mathrm{LT}(\mathscr{S}^1(f_1, f_2) \right\rangle = aY \left\langle X, \frac{b}{a}Y, b \right\rangle \\ & \subsetneq \left\langle \mathrm{LT}(\mathscr{S}^2(f_1, f_2) \right\rangle = \left\langle aXY, bY^2, abY, a^2b \right\rangle = \mathrm{MLT}(\left\langle f_1, f_2 \right\rangle) \end{split}$$

Corollary 4.15 (Discrete coherent context 2.2). Let $I = \langle f_1, \ldots, f_p \rangle$ be a nonzero finitely generated submodule of \mathbf{H}_n^m . Suppose that $\mathrm{MLT}(I)$ is finitely generated, i.e. that there exist $u_1, \ldots, u_t \in I$ such that $\mathrm{MLT}(I) = \langle \mathrm{LT}(u_1, \ldots, u_t) \rangle$. Then there exists $q \in \mathbb{N}$ such that $\mathrm{MLT}(I) = \langle \mathrm{LT}(\mathscr{S}^q(f_1, \ldots, f_p)) \rangle$.

Remark 4.16. In Corollary 4.15 with m = 1, if the considered monomial order refines total degree, then, writing $u_k = \sum_{j=1}^p g_{k,j} f_j$ with $g_{k,j} \in \mathbf{R}[\underline{X}]$ and letting $d = \max_{1 \le k \le t, 1 \le j \le p} (\operatorname{tdeg}(g_{k,j}) + \operatorname{tdeg}(f_j))$ and $\delta = \min_{1 \le k \le t} \operatorname{tdeg}(u_k)$ (assumed to be ≥ 1), we have $q \le \binom{n+d}{d} - \binom{n+\delta-1}{\delta-1}$.

5 Basic algorithms

The division algorithm

This algorithm in Simple division context 2.3 needs **R** to be strongly discrete; note that coherence is not used here. Like the classical division algorithm for $\mathbf{F}[\underline{X}]^m$ with **F** a discrete field (see Yengui 2015, Algorithm 211), this algorithm has the following

goal.

Definition 5.1. The vector r is called a remainder of u on division by the list $H = h_1, \ldots, h_p$ and is denoted by $r = \overline{u}^H$.

Algorithm 5.2 provides a suitable answer: a suitable remainder r and suitable quotients q_i . Nevertheless, there are a priori many different possible answers.

Division algorithm 5.2 (Simple division context 2.3).

```
<sup>1</sup> Division (u, h_1, \ldots, h_p)
   local variables j: [1 \dots p], D: subset of [1 \dots p],
 2
                                   c, c_1, \ldots, c_n, a_1, \ldots, a_n : \mathbf{R}, M, M_1, \ldots, M_n : \mathbf{H}_n^m;
 з
 r \leftarrow 0;
   for j from 1 to p do q_j \leftarrow 0; M_j \leftarrow LM(h_j); c_i \leftarrow LC(h_j) od;
    while u \neq 0 do
       M \leftarrow \operatorname{LM}(u); \ c \leftarrow \operatorname{LC}(u); \ D \leftarrow \{j; M_i \mid M\};
 7
        if c \in \langle c_j; j \in D \rangle then
 8
            find (a_j)_{j\in D} such that \sum_{j\in D} a_j c_j = c;
9
           u \leftarrow u - \sum_{j \in D} a_j (M/M_j) h_j;
10
            for j \in D do q_j \leftarrow q_j + a_j(M/M_j) od
11
        else r \leftarrow r + cM; u \leftarrow u - cM fi
12
13 od :
14 return r, q_1, \ldots, q_p
```

One checks by induction that $LM(q_j) LM(h_j) \le LM(u)$ and $u = q_1h_1 + \cdots + q_ph_p + r$.

Syzygy algorithms

Notation 5.3. We denote by List(A) the set of (finite) lists of elements of A.

These algorithms take place in Discrete coherent context 2.2. They are a key tool for constructing a Gröbner basis and have been introduced by Buchberger (1965) for the case where the base ring is a discrete field.

We begin with the basic syzygy algorithm returning $S^E(a_1M_1,\ldots,a_pM_p) = S_1^E,\ldots,S_\ell^E$ for $a_1M_1,\ldots,a_pM_p \in \mathbf{H}_n^m$ and a subset $E \subseteq \mathscr{P}(M_1,\ldots,M_p)$: see Defini-

tion 4.1. Let us recall that $\mathbf{H}_n^1 = \mathbf{R}[\underline{X}]$.

Input a_1M_1, \ldots, a_pM_p terms in $\mathbf{H}_n^m, E \subseteq \mathcal{P}(M_1, \ldots, M_p)$. Output A list of syzygies $(S_{1,j}^E)_{j \in [1..p]}, \ldots, (S_{\ell,j}^E)_{j \in [1..p]}$ for (a_1M_1, \ldots, a_pM_p) such that $S_{i,j}^E = 0$ for $j \notin E$ and, for every syzygy expression $g_1a_1M_1 + \cdots + g_pa_pM_p$ whose summands have leading monomial index set E, $(\mathrm{LT}(g_j))_{j \in E} \in \langle (S_{1,j}^E)_{j \in E}, \ldots, (S_{\ell,j}^E)_{j \in E} \rangle$.

Basic syzygy algorithm for terms 5.4 (basic syzygies of terms, Definition 4.1, Discrete coherent context 2.2).

1 BasicSyzygiesOfTerms $(a_1M_1, \ldots, a_pM_p, E)$ 2 local variables $j : [1 \dots p]$, $\ell, i : \mathbb{N}$, $s_1, \ldots, s_\ell : \mathbb{R}^E$, $M^E : \mathbb{H}_n^m$; 3 find $\ell, s_1, \ldots, s_\ell$ such that $\operatorname{Syz}((a_j)_{j \in E}) = \langle s_1, \ldots, s_\ell \rangle$; 4 for i from 1 to ℓ do 5 for j from 1 to p do 6 $M^E \leftarrow \operatorname{lcm}(M_j; j \in E);$ 7 if $j \in E$ then $S_{i,j}^E \leftarrow s_{i,j}(M^E/M_j)$ else $S_{i,j}^E \leftarrow 0$ fi 8 od 9 od; 10 return $(S_{1,j}^E)_{j \in [1..p]}, \ldots, (S_{\ell,j}^E)_{j \in [1..p]}$

We now give an algorithm whose goal is to provide a generating set of syzygies for a vector of terms in \mathbf{H}_n^m ; see Definition 4.1 and Proposition 4.4.

Input $a_1M_1 = a_1M'_1e_{i_1}, \ldots, a_pM_p = a_pM'_pe_{i_p}$ terms in \mathbf{H}_n^m . Output a list of syzygies $S_i^E \in \mathbf{R}[\underline{X}]^p$ such that the S_i^E 's generate $\operatorname{Syz}(a_1M_1, \ldots, a_pM_p)$.

In the algorithm, we construct the syzygies S_i^E by successive concatenations of the lists obtained by the previous algorithm.

Syzygy algorithm for terms 5.5 (syzygies of terms, see Definition 4.1, Discrete coherent context 2.2).

```
1 SyzygiesOfTerms(a_1M_1, \ldots, a_pM_p)

2 local variables E: subset of [1 \dots p], S^E: List(\mathbf{R}[\underline{X}]^p);

3 S \leftarrow;

4 for E in \mathcal{P}(M_1, \dots, M_p) do

5 S^E \leftarrow BasicSyzygiesOfTerms(a_1M_1, \dots, a_pM_p, E);

6 S \leftarrow S, S^E

7 od;

8 return S
```

In the case of an ideal (m = 1), one may forget about the basis vectors e_{i_1}, \ldots, e_{i_p} and one has $\mathscr{P}(M_1, \ldots, M_p) = \mathscr{P}_p$.

S-list algorithms

We have the following goal corresponding to the S-list $\mathscr{S}(f_1, \ldots, f_p)$ in Definition 4.9 (Discrete coherent context 2.2).

Input $f_1, \ldots, f_p \in \mathbf{H}_n^m$ not all zero, **Output** The S-list $\mathscr{S} = \mathscr{S}(f_1, \ldots, f_p)$ of f_1, \ldots, f_p as in Definition 4.9.

S-list algorithm 5.6 (S-list algorithm, Definition 4.9, Discrete coherent context 2.2).

```
1 Slist (f_1, \ldots, f_p)

2 local variables Si, Ss : \mathbf{R}[\underline{X}]^p; S, Slist : \text{List}(\mathbf{R}[\underline{X}]^p);

3 \mathcal{S} \leftarrow;

4 S \leftarrow SyzygiesOfTerms(LT(f_1, \ldots, f_p));

5 for Si in S do

6 Ss \leftarrow Si_1f_1 + \cdots + Si_pf_p;

7 if Ss \neq 0 then \mathcal{S} \leftarrow \mathcal{S}, Ss fi

8 od;

9 return \mathcal{S}
```

We have the following goal corresponding to the S-list $\mathscr{S}^q(f_1,\ldots,f_p)$ in Definition 4.9 (Discrete coherent context 2.2.)

Input $q \in \mathbb{N}, f_1, \ldots, f_p \in \mathbf{H}_n^m$ not all zero, **Output** The iterated S-list $\mathscr{S}^q = \mathscr{S}^q(f_1, \ldots, f_p)$ of f_1, \ldots, f_p as in Definition 4.9.

S-list algorithm 5.7 (iterated S-list algorithm, Definition 4.9, Discrete coherent context 2.2).

```
1 IteratedSlist (q, f_1, ..., f_p)

2 local variables r : \mathbb{N};

3 \mathscr{S}^q \leftarrow f_1, ..., f_p;

4 for r from 1 to q do \mathscr{S}^q \leftarrow \mathscr{S}^q,Slist(\mathscr{S}^q) od;

5 return \mathscr{S}^q
```

Rewriting algorithms

The next algorithm corresponds to Proposition 4.6.

Input $f_1, \ldots, f_p \in \mathbf{H}_n^m$ not all zero; $g_1, \ldots, g_p \in \mathbf{R}[\underline{X}]$; the leading monomial L of the summands of $\sum_{j=1}^p g_j f_j = u$ (let E be the corresponding leading monomial index set) is > LT(u).

Output $f_1, \ldots, f_{p+\ell}$ in $\mathscr{S}(f_1, \ldots, f_p)$ extending $f_1, \ldots, f_p; g_1, \ldots, g_{p+\ell} \in \mathbf{R}[\underline{X}]$ with the original g_j 's replaced by $g_j - \mathrm{LT}(g_j)$ for $j \in E$ and unchanged outside E; $u = \sum_{j=1}^{p+\ell} g_j f_j$ and the leading monomial of its summands is < L.

Rewriting algorithm 5.8 (rewriting a linear combination, Proposition 4.6, Discrete coherent context 2.2).

Rewriting $((g_1, f_1), \ldots, (g_p, f_p))$ 1 local variables $j: [1 \dots p]$, $a_1, \dots, a_p, b_1, \dots, b_p, c_1, \dots, c_\ell : \mathbf{R}$, 2 $M_1, \ldots, M_p, L, M^E : \mathbf{H}_n^m, N_1, \ldots, N_p : \mathbf{R}[X],$ 3 $\ell, i: \mathbb{N}, E:$ subset of $[1 \dots p],$ 4 $(S_{1,j}^E)_{j\in[1..p]}, \ldots, (S_{\ell,j}^E)_{j\in[1..p]} : \mathbf{R}[\underline{X}]^p, s_1, \ldots, s_\ell : \mathbf{R}^E;$ 5 6 for j in $[1 \dots p]$ do $a_i \leftarrow \mathrm{LC}(f_i); \ M_i \leftarrow \mathrm{LM}(f_i); \ b_i \leftarrow \mathrm{LC}(g_i); \ N_i \leftarrow \mathrm{LM}(g_i) \ \mathsf{od};$ 7 $L \leftarrow \sup\{N_i M_j; j \in [1 ... p]\}; E \leftarrow \{j \in [1 ... p]; N_j M_j = L\};$ 9 $M^E \leftarrow \operatorname{lcm}(M_j; j \in E);$ ${}_{^{10}} (S^E_{1,j})_{j \in [1..p]}, \dots, (S^E_{\ell,j})_{j \in [1..p]} \leftarrow \mathsf{BasicSyzygiesOfTerms}(a_1M_1, \dots, a_pM_p, E);$ 11 for i in $[1..\ell]$ do $s_i \leftarrow (\operatorname{LC}(S_{i,i}^E))_{i \in E}$ od; 12 find $(c_i)_{i \in [1..\ell]}$ such that $(b_j)_{j \in E} = \sum_{i \in [1..\ell]} c_i s_i$; 13 for i in $[1 \dots \ell]$ do $g_{p+i} \leftarrow c_i \frac{L}{M^E}$; $f_{p+i} \leftarrow \sum_{j \in E} S^E_{i,j} f_j$ od; 14 for j in E do $g_j \leftarrow g_j - LT(g_j)$ od; ¹⁵ return $(g_1, f_1), \ldots, (g_{p+\ell}, f_{p+\ell})$

We have the following goal corresponding to Fundamental theorem 4.11 in Discrete coherent context 2.2. This is an iteration of the previous one with a counter q for the number of iterations.

Input $f_1, \ldots, f_p \in \mathbf{H}_n^m \setminus \{0\}, g_1, \ldots, g_p \in \mathbf{R}[\underline{X}]$; we let $u = \sum_{j=1}^p g_j f_j$. Output $q \in \mathbb{N}$ and f_1, \ldots, f_t in $\mathscr{S}^q(f_1, \ldots, f_p)$ extending f_1, \ldots, f_p such that $\mathrm{LT}(u) \in \langle \mathrm{LT}(f_1, \ldots, f_t) \rangle$.

Rewriting algorithm 5.9 (iterated rewriting of a linear combination, Fundamental theorem 4.11, Discrete coherent context 2.2).

1 Iterated Rewriting ($(g_1, f_1), \ldots, (g_p, f_p)$) 2 $q \leftarrow 0$; ³ while $LM(g_1f_1 + \dots + g_pf_p) < \sup\{LM(g_1) LM(f_1), \dots, LM(g_p) LM(f_p)\}$ do ⁴ $(g_1, f_1), \dots, (g_{p+\ell}, f_{p+\ell}) \leftarrow \mathsf{Rewriting}((g_1, f_1), \dots, (g_p, f_p));$ ⁵ $p \leftarrow p + \ell; q \leftarrow q + 1$ od; ⁶ return q, f_1, \dots, f_p

6 Buchberger's algorithm

The proof of the following theorem parallels exactly the proof of the analogue Theorem 4.2.3 of Adams and Loustaunau (1994).

Fundamental theorem 6.1 (Strongly discrete coherent context 2.5).

(1) Buchberger's criterion. Let $f_1, \ldots, f_p \in \mathbf{H}_n^m$ not all zero, and denote by S_1, \ldots, S_ℓ the generators of $\operatorname{Syz}(\operatorname{LT}(f_1, \ldots, f_p))$ computed in Proposition 4.4. Then $G = f_1, \ldots, f_p$ is a Gröbner basis for $\langle f_1, \ldots, f_p \rangle$ if and only if for every $1 \leq i \leq \ell$, we have

$$\overline{S_{i,1}f_1 + \dots + S_{i,p}f_p}^G = 0.$$

(2) Buchberger's algorithm works. Let $f_1, \ldots, f_p \in \mathbf{H}_n^m \setminus \{0\}$ and $M = \langle f_1, \ldots, f_p \rangle$. If the module of leading terms MLT(M) of the module M is finitely generated, then the (generalised) Buchberger algorithm 6.2 computes a Gröbner basis for $\langle f_1, \ldots, f_p \rangle$.

The (generalised) Buchberger algorithm has the following goal.

Input $f_1, \ldots, f_p \in \mathbf{H}_n^m \setminus \{0\}$. Output a Gröbner basis $f_1, \ldots, f_p, \ldots, f_t$ for $\langle f_1, \ldots, f_p \rangle$ extending f_1, \ldots, f_p .

Buchberger's algorithm 6.2 (Strongly discrete coherent context 2.5).

```
<sup>1</sup> Buchberger(f_1, \ldots, f_p)
 <sup>2</sup> local variables \mathscr{G}: List(\mathbf{H}_n^m); f, r: \mathbf{H}_n^m; L: List(\mathbf{R}[\underline{X}]);
 G \leftarrow f_1, \ldots, f_p;
    repeat
 4
         \mathscr{S} \leftarrow \mathsf{Slist}(G);
 5
         for f in \mathcal{S} do
 6
             remove f from \mathcal{S};
 7
             r, L \leftarrow Division (f, G);
 8
             if r \neq 0 then \mathscr{G} \leftarrow \mathscr{G}, r fi
 9
         od :
10
         if \mathscr{S} \neq \emptyset then G \leftarrow G, \mathscr{S} fi
11
     until \mathcal{S} = \emptyset;
12
13 return G
```

Remark 6.3. If the algorithm terminates, then we can transform the obtained Gröbner basis into a Gröbner basis $h_1, \ldots, h_{p'}$ such that no term of an element h_{ℓ} lies in $\langle \operatorname{LT}(h_k) ; k \neq \ell \rangle$ by replacing each element of the Gröbner basis with a remainder of it on division by the other nonzero elements and by repeating this process until it stabilises. Such a Gröbner basis is called a *pseudo-reduced* Gröbner basis. The terminology of "reduced" Gröbner basis is used only in the case where a way of normalising its elements is specified: see Yengui 2015, Remark 239.

7 Schreyer's syzygy algorithm

Definition 7.1 (Schreyer's monomial order). Let **R** be a discrete ring. Consider a list $G = f_1, \ldots, f_p$ in $\mathbf{H}_n^m \setminus \{0\}$ and the finitely generated submodule $U = \langle G \rangle = \mathbf{R}[\underline{X}]f_1 + \cdots + \mathbf{R}[\underline{X}]f_p$ of \mathbf{H}_n^m . Let $(\epsilon_1, \ldots, \epsilon_p)$ be the canonical basis of $\mathbf{R}[\underline{X}]^p$. Schreyer's monomial order induced by > and G on $\mathbf{R}[\underline{X}]^p$ is the order $>_G$ defined as follows:

$$\underline{X}^{\alpha} \epsilon_k >_G \underline{X}^{\beta} \epsilon_j \quad \text{if} \quad \begin{vmatrix} \text{either } \operatorname{LM}(\underline{X}^{\alpha} f_k) > \operatorname{LM}(\underline{X}^{\beta} f_j) \\ \text{or both } \operatorname{LM}(\underline{X}^{\alpha} f_k) = \operatorname{LM}(\underline{X}^{\beta} f_j) \text{ and } k < j. \end{vmatrix}$$

Schreyer's monomial order is defined on $\mathbf{R}[\underline{X}]^p$ in the same way as when \mathbf{R} is a discrete field (see Ene and Herzog 2012, p. 66). Note that it actually depends only on the leading monomials $\mathrm{LM}(G)$ of G.

Now we shall follow closely the ingenious proof by Schreyer (1980) of Hilbert's syzygy theorem via Gröbner bases, but with a strongly discrete coherent ring instead of a field. Schreyer's proof is very well explained in Ene and Herzog 2012, §§ 4.4.1–4.4.3.

Schreyer's syzygy algorithm below takes also place in Strongly discrete coherent context 2.5 for \mathbf{R} . It has the following goal.

Input a Gröbner basis f_1, \ldots, f_p for a submodule of \mathbf{H}_n^m . **Output** a Gröbner basis $(u_i^E)_{1 \le i \le \ell^E, E \in \mathscr{P}(\mathrm{LM}(f_1), \ldots, \mathrm{LM}(f_p))}$ for $\mathrm{Syz}(f_1, \ldots, f_p)$ with respect to Schreyer's monomial order induced by > and f_1, \ldots, f_p .

Schreyer's syzygy algorithm 7.2 (Strongly discrete coherent context 2.5).

```
\begin{array}{ll} & \operatorname{SchreyerSyzygy}\left(f_{1},\ldots,f_{p}\right)\\ & 2 \ \operatorname{local} \ \operatorname{variables} \ \left(S_{i}^{E}\right)_{1\leq i\leq \ell^{E},E\in\mathscr{P}(\operatorname{LM}(f_{1}),\ldots,\operatorname{LM}(f_{p}))}:\operatorname{List}(\mathbf{R}[\underline{X}]^{p}),\\ & 3 \ & q_{1},\ldots,q_{p}:\mathbf{R}[\underline{X}];\\ & 4 \ \left(S_{i}^{E}\right)_{1\leq i\leq \ell^{E},E\in\mathscr{P}(\operatorname{LM}(f_{1}),\ldots,\operatorname{LM}(f_{p}))}\leftarrow\operatorname{SyzygiesOfTerms}(f_{1},\ldots,f_{p});\\ & 5 \ \operatorname{for} \ E \ \operatorname{in} \ \mathscr{P}(\operatorname{LM}(f_{1}),\ldots,\operatorname{LM}(f_{p})) \ \operatorname{do}\\ & 6 \ & \text{for} \ i \ \operatorname{from} \ 1 \ \operatorname{to} \ l^{E} \ \operatorname{do}\\ & 7 \ & 0,q_{1},\ldots,q_{p}\leftarrow\operatorname{Division}\left(S_{i,1}^{E}f_{1}+\cdots+S_{i,p}^{E}f_{p},f_{1},\ldots,f_{p}\right)\\ & 8 \ & (\operatorname{note} \ \operatorname{that} \operatorname{LM}(S_{i,1}^{E}f_{1}+\cdots+S_{i,p}^{E}f_{p})\geq\operatorname{LM}(q_{j}f_{j}) \ \operatorname{whenever} q_{j}f_{j}\neq 0); \end{array}
```

9 $u_i^E \leftarrow S_{i,1}^E \epsilon_1 + \dots + S_{i,p}^E \epsilon_p - q_1 \epsilon_1 - \dots - q_p \epsilon_p$ 10 od
11 od;
12 return $(u_i^E)_{1 \leq i \leq \ell^E, E \in \mathscr{P}(\mathrm{LM}(f_1), \dots, \mathrm{LM}(f_p))}$

Note that the polynomials q_1, \ldots, q_p of lines 7–8 may have been computed while constructing the Gröbner basis.

Remark 7.3. For an arbitrary system of generators h_1, \ldots, h_p for a submodule U of \mathbf{H}_n^m , the syzygy module of h_1, \ldots, h_p is easily obtained from the syzygy module of a Gröbner basis for U (see Yengui 2015, Theorem 296).

Fundamental theorem 7.4 (Schreyer's algorithm, Strongly discrete coherent context 2.5). Let U be a submodule of \mathbf{H}_n^m with Gröbner basis f_1, \ldots, f_p . Then the relations u_i^E computed by Schreyer's syzygy algorithm 7.2 form a Gröbner basis for the syzygy module $\operatorname{Syz}(f_1, \ldots, f_p)$ with respect to Schreyer's monomial order induced by > and f_1, \ldots, f_p . Moreover, for E a position level subset of $[1 \ldots p]$ and $1 \le i \le \ell^E$,

$$\operatorname{LT}(u_i^E) = s_{i,r}^E \ M^E / M_r \ \epsilon_r \ \text{with} \ r = \min\{j \in E \ ; \ s_{i,j}^E \neq 0\}$$

$$\tag{4}$$

in the notation of Definition 4.1 with $M_1 = LM(f_1), \ldots, M_p = LM(f_p)$.

Proof (a slight modification of the proof of Ene and Herzog 2012, Theorem 4.16). Let us use the notation of Schreyer's syzygy algorithm 7.2. Recall that $u_i^E = (S_{i,1}^E - q_1)\epsilon_1 + \dots + (S_{i,p}^E - q_p)\epsilon_p$, $S_{i,1}^E f_1 + \dots + S_{i,p}^E f_p = q_1f_1 + \dots + q_pf_p$, and $\mathrm{LM}(q_jf_j) \leq \mathrm{LM}(S_{i,1}^E f_1 + \dots + S_{i,p}^E f_p) < (M^E/M_k)M_k = M^E$ for any $k \in E$. So $\mathrm{LT}(u_i^E) = \mathrm{LT}(S_i^E) = s_{i,r}^E M^E/M_r \epsilon_r$ where $r = \min\{j \in E : s_{i,j}^E \neq 0\}$.

Let us show now that the relations u_i^E form a Gröbner basis for the syzygy module $\operatorname{Syz}(f_1,\ldots,f_p)$. For this, let $v = v_1\epsilon_1 + \cdots + v_p\epsilon_p \in \operatorname{Syz}(f_1,\ldots,f_p)$ and let us show that $\operatorname{LT}(v) \in \langle \operatorname{LT}(u_i^E) ; 1 \leq i \leq \ell^E, E \in \mathscr{P}(M_1,\ldots,M_p) \rangle$. Let us write $\operatorname{LM}(v_j\epsilon_j) = N_j\epsilon_j$ and $\operatorname{LC}(v_j\epsilon_j) = c_j$ for $1 \leq j \leq p$. Then $\operatorname{LM}(v) = N_k\epsilon_k$ for some $1 \leq k \leq p$. Now let $v' = \sum_{j \in D} c_j N_j\epsilon_j$, where D is the set of those j for which $N_jM_j = N_kM_k$. By definition of Schreyer's monomial order, we have $j \geq k$ for all $j \in D$. Substituting each ϵ_j in v' by $T_j = \operatorname{LT}(f_j)$, the sum becomes zero. Therefore v' is a syzygy of the terms T_j with $j \in D$. By virtue of Proposition 4.4, v' is a linear combination of elements in $S((T_j)_{j\in D})$ of the form S_i^E with $E \subseteq D$ and $1 \leq i \leq \ell^E$. By inspecting the jth component of v', we deduce that there exist $w_1,\ldots,w_t \in \mathbf{R}[\underline{X}]$, position level subsets E_1,\ldots,E_t of D with $j \in E_1 \cap \cdots \cap E_t$, nonnegative integers $1 \leq i_1 \leq \ell_{E_1},\ldots,1 \leq i_t \leq \ell_{E_t}$, such that $c_jN_j = w_1s_{E_1,i_1,j}M_{E_1}/M_j + \cdots + w_ts_{E_t,i_t,j}M_{E_t}/M_j$, and $s_{E_1,i_1,j},\ldots,s_{E_t,i_t,j} \neq 0$. As j > k for all $j \in D \setminus \{k\}$, it follows that $\operatorname{LT}(v') \in \langle \operatorname{LT}(S_{E_1,i_1},\ldots,S_{E_t,i_t})\rangle$. The desired result follows since $\operatorname{LT}(v) = \operatorname{LT}(v')$.

Schreyer's monomial order is a tailor-made term over position monomial order which changes at each iteration, i.e. after each computation of a Gröbner basis of the syzygy module of the considered Gröbner basis. Schreyer's trick is, for $v = v_1 \epsilon_1 + \cdots + v_p \epsilon_p \in \operatorname{Syz}(f_1, \ldots, f_p)$, to prioritise (by deciding that they are greater) the $\operatorname{LM}(v_j \epsilon_j)$ such that $\operatorname{LM}(v_j f_j) = \max(\operatorname{LM}(v_1 f_1), \ldots, \operatorname{LM}(v_p f_p))$, and to order the obtained generators u_i^E of $\operatorname{Syz}(f_1, \ldots, f_p)$ in such a way that X_n does not appear in the leading terms of the u_i^E (when computing a Gröbner basis for the first syzygy module), and to iterate this process until exhausting all the indeterminates $X_n, X_{n-1}, \ldots, X_1$ from the leading terms of the Gröbner basis of the syzygy module. Once we reach this situation, we continue the resolution over the base ring \mathbf{R} .

As a consequence of Theorem 7.4, we obtain the following constructive version of Hilbert's syzygy theorem for a strongly discrete coherent ring.

Theorem 7.5 (Hilbert's syzygy theorem, Strongly discrete coherent context 2.5). Let \mathbf{H}_n^m be a free $\mathbf{R}[\underline{X}]$ -module with basis (e_1, \ldots, e_m) , and > a monomial order on \mathbf{H}_n^m . Let U be a finitely generated submodule of \mathbf{H}_n^m such that MLT(U) is finitely generated with respect to some monomial order. Then $M = \mathbf{H}_n^m/U$ admits an $\mathbf{R}[\underline{X}]$ -resolution

$$0 \to F_q/V \to F_{q-1} \to \cdots \to F_1 \to F_0 \to M \to 0$$

such that $q \leq n+1, F_0, \ldots, F_q$ are finitely generated free $\mathbf{R}[\underline{X}]$ -modules, and V is generated by finitely many iterated syzygies whose leading terms with respect to Schreyer's induced monomial order do not depend on the indeterminates X_n, \ldots, X_1 .

Proof. Let (f_1, \ldots, f_p) be a Gröbner basis for U with respect to the considered order. Reorder the f_j 's so that whenever $\operatorname{LM}(f_j)$ and $\operatorname{LM}(f_k)$ involve the same position for some k < j, say $\operatorname{LM}(f_k) = M'_k e_{i_k}$ and $\operatorname{LM}(f_j) = M'_j e_{i_j}$ with $i_j = i_k$, then $\operatorname{deg}_{X_n}(M'_k) \ge \operatorname{deg}_{X_n}(M'_j)$. Consider the leading monomials $\operatorname{LM}(u^E_i) = M^E/M_r \epsilon_r$, $r = \min\{j \in E ; s^E_{i,j} \neq 0\}$, computed by Schreyer's syzygy algorithm 7.2: our reordering entails that $\operatorname{deg}_{X_n} M_E = \operatorname{deg}_{X_n} M_r$, so that the indeterminate X_n does not appear in $\operatorname{LM}(u^E_i)$. Thus, after at most n computations of the iterated syzygies, we reach the desired situation.

Remark 7.6. This theorem generalises Theorems 5.5 and 6.2 of Gamanda, Lombardi, Neuwirth, and Yengui 2020. Note that their statement there, as well as that of its Theorem 5.9, needs to be amended: "the TOP lexicographic monomial order" needs to be replaced by "some monomial order"; the proof given here shows how their proof needs to be amended. See the arXiv version Gamanda, Lombardi, Neuwirth, and Yengui 2024.

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